

A technique for solving stochastic programs with first stage integer variables

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Outline

1. Review of current solution approaches for Integer SP.
2. Discussion of our approach.
3. Analysis of Bounds.
4. Preliminary computational results.



New Perspective of Investment Decisions

The issues

DCF is inadequate...

Three leading characteristics

- ◆ Investment Decisions (costs) irreversible
- ◆ Future returns are uncertain
- ◆ Another key aspect is timing
 - ◆ Invest
 - ◆ Disinvest
 - ◆ Not invest...postpone

Are all strategic decisions



Decision making under uncertainty

Buy flexibility

≡

Hedge against uncertainty

≡

Make robust decisions



General two-stage SP

$$\text{Min } c^T x + \int p(\omega) f(\omega) y(\omega)$$

Subject to

$$\begin{aligned} Ax &= b, \\ B(\omega)x + D(\omega)y(\omega) &= h(\omega) \\ x \geq 0, y &\geq 0. \end{aligned}$$

ω = random event,

$p(\omega)$ = probability,

$f(\omega)$ = Second - stage objective,

$B(\omega)$ = Technical matrix,

$D(\omega)$ = Recourse matrix,

$h(\omega)$ = Right - hand side,

x = First - stage decision ,

$y(\omega)$ = Second - stage decision

let $\xi_\omega = \{f, h, B, D\}_\omega$

$$\text{Min } c^T x + E_\xi Q(x, \xi)$$

Subject to

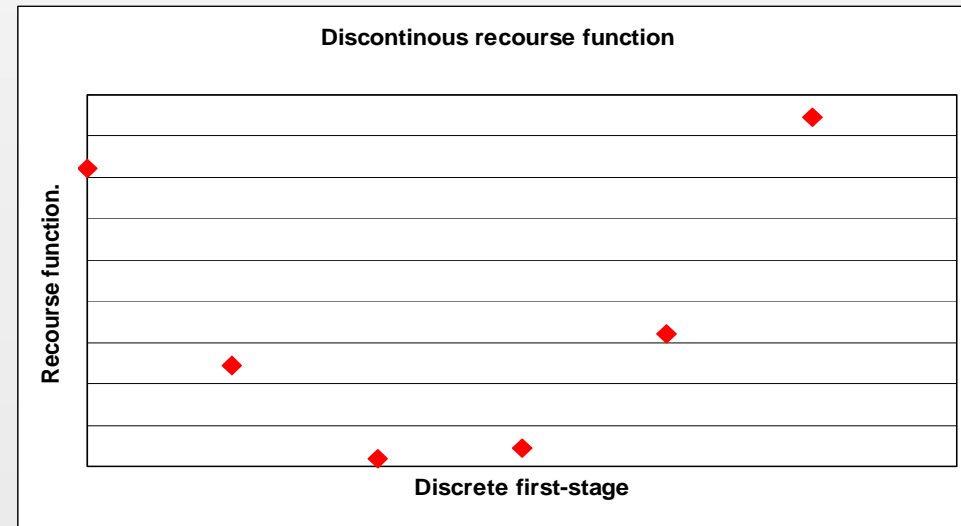
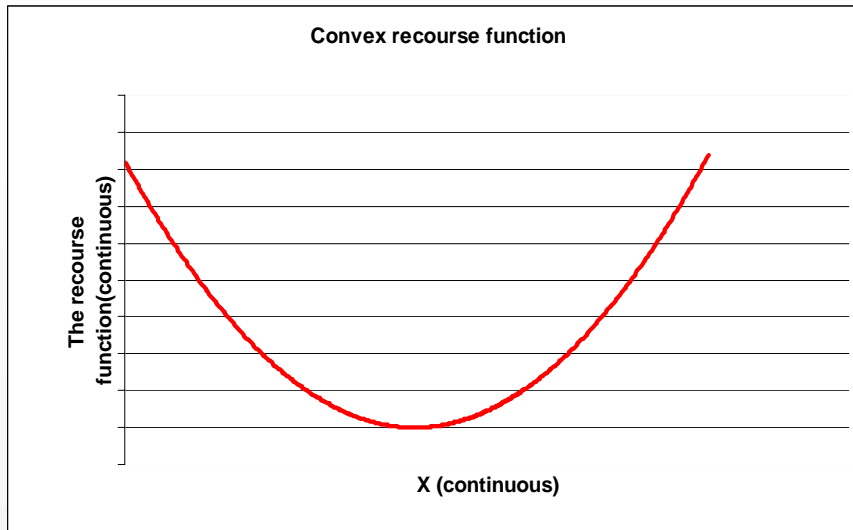
$$Ax = b,$$

where

$$Q(x, \xi) = \text{Min} \{ f y(\omega) \mid D y(\omega) = h - B x, y(\omega) \geq 0 \}$$



Shape of the recourse function



Stochastic integer programs

$x \in Z^{n_1}$, OR $y(\omega) \in Z^{n_2}$, OR $x \in Z^{n_1}$ AND $y(\omega) \in Z^{n_2}$

Applications

First Stage integer

- Supply chain planning, Mitra et al.
- Asset allocation, Mulvey and Rusczyński
- Electric utility planning, Bienstock and Shapiro
- Goods distribution, Cheung and Powell

Second Stage integer

- Scheduling decisions, Dempster et al.
- Routing decisions Spaccamela et al.
- Fixed-charge, Bitran et al.
- Change over costs, Caroe et al.



Current approaches

- Cutting plane - *Wollmer*
- Scenario decomposition- *Rockafellar and Wets*
- Integer L-shaped – *Laporte and Louveaux*
- Convex hull for Simple integer recourse – *Haneveld, Stougie, Van der Vlerk*



Current approaches

- Progressive hedging - *Takriti et al.*
- Groebner bases – *Schultz et al.*
- Integer L-shaped – *Caroe and Tind*
- Dual decomposition – *Caroe and Schultz*
- Branch and bound based approach – *Ahmed et al.*



Problem Statement

$$\text{Min } c^T x + E_{\xi} Q(x, \xi)$$

Subject to

$$Ax = b,$$

where

$$Q(x, \xi) = \text{Min} \{ fy(\omega) \mid Dy(\omega) = h - Bx \}$$

$$x = (x', x''), x \in \{0,1\}^{n_1}, x^{n_1} \in \mathcal{R}^{n_1}, y(\omega) \in \mathcal{R}^{n_2}$$



Motivation for the algorithm

A Strategic Supply Chain model

Rows	6768
Discrete	2096
Continuous	54400
Non-Zeroes	1154034
Scenarios	100



Assumptions

- Relatively complete recourse
- Discrete distribution



The Approach

1. Lagrangean relaxation used to generate a candidate set of integer feasible solutions for each WS model.
2. Construct a convex hull of this set of integer feasible solutions.
3. Not necessarily a convex hull to the original 2-stage ISP.
4. We compute bounds to the 2-stage SP model and qualify the approximated model.



Algorithm

P_{2SP}

$$\min Z = cx + \sum_s p_s f_s y_s$$

subject to

$$Ax = b$$

$$B_s x + D_s y_s \leq d_s \quad \forall s \in \{1, \dots, S\}$$

$$x \in \{0,1\}^{n_1}$$

$$y_s \geq 0$$

$P_{WS}(s)$

$$\min Z = cx + f_s y_s$$

subject to

$$Ax = b$$

$$B_s x + D_s y_s \leq d_s$$

$$x \in \{0,1\}^{n_1}$$

$$y_s \geq 0$$

Discrete Variables

Continuous Variables



Algorithm

$P_{WS}(s)$

$$\begin{aligned} \min Z_{WS} &= cx + f_s y_s \\ \text{subject to} \\ Ax &= b \\ B_s^R x + D_s^R y_s &\leq d_s^R \\ B_s^{NR} x + D_s^{NR} y_s &\leq d_s^{NR} \\ x \in \{0,1\}^{n_1}, y &\geq 0 \end{aligned}$$

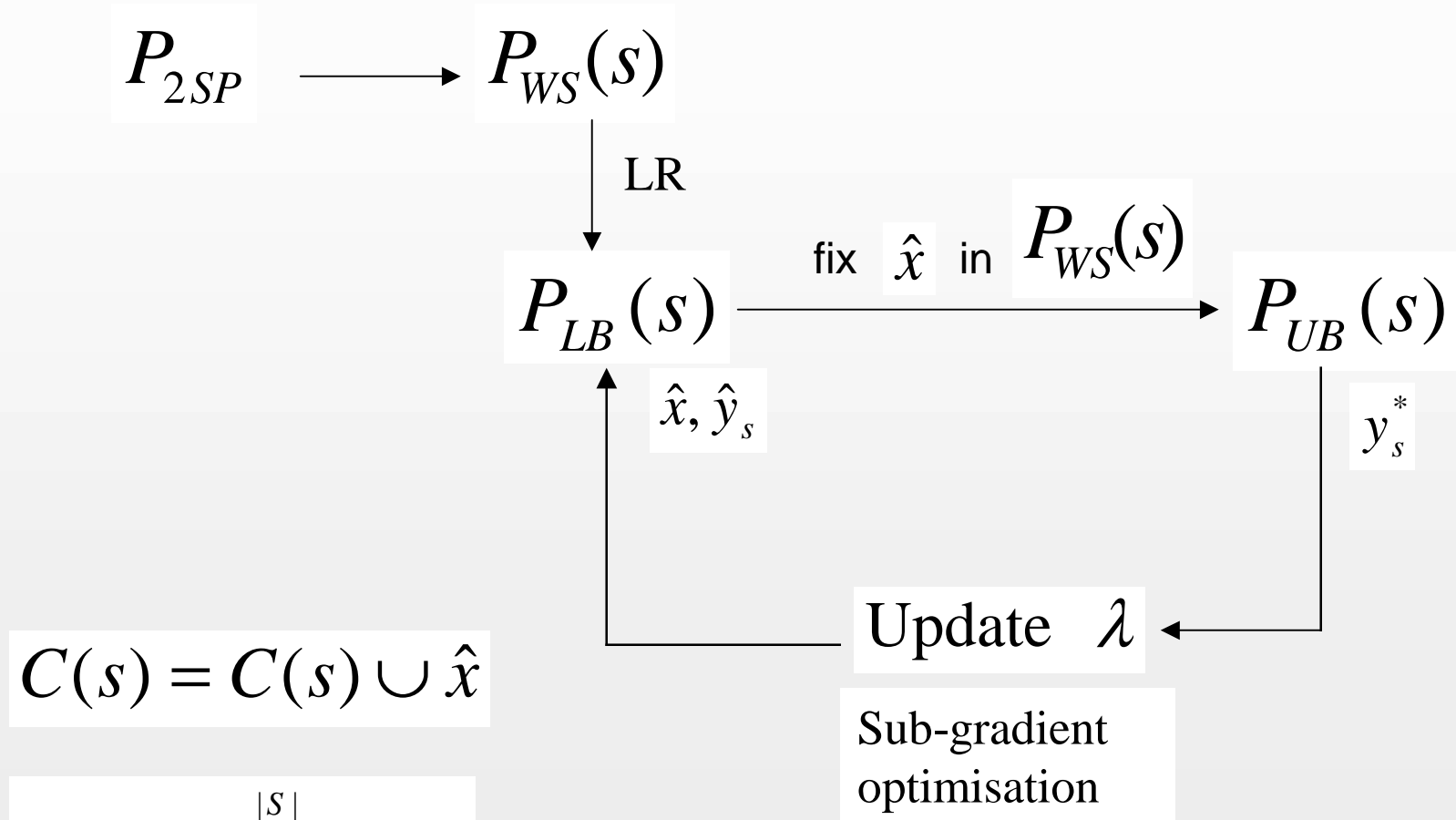
$P_{LB}(s)$

$$\begin{aligned} \min Z_{LB} &= (c + \lambda B_s^R)x + (f_s + D_s^R)y_s^R - \lambda_s^R d_s^R \\ \text{subject to} \\ Ax &= b \\ B_s^{NR} x + D_s^{NR} y_s^{NR} &\leq d_s^{NR} \\ x \in \{0,1\}^{n_1}, y &\geq 0 \end{aligned}$$

$P_{UB}(s)$

$$\begin{aligned} \min Z_{UB} &= c\hat{x} + f_s y_s \\ \text{subject to} \\ A\hat{x} &= b \\ B_s \hat{x} + D_s y_s &\leq d_s \\ \hat{x} \in \{0,1\}^{n_1}, y &\geq 0 \end{aligned}$$

Algorithm



$$C(s) = C(s) \cup \hat{x}$$

$$ACS = \bigcup_{s=1}^{|S|} C(s)$$



Algorithm

Stopping Criteria:

1. $\frac{\text{BUB} - \text{BLB}}{|\text{BLB}| + 1} < \text{Tolerance}$
2. Pass \geq Maximum Number of Iterations
3. Satisfaction of the relaxed constraints, $B_s^R \hat{x} + D_s^R \hat{y}_s \leq d_s^R$



The Approximated model

APXP_{2SP}

$$\begin{aligned} \min Z &= cx + \sum_s p_s fy_s \\ \text{subject to} \\ Dy_s &\leq h - Bx, \\ Ey_s &= d_s, \quad \forall s \in \{1, \dots, S\} \\ x &= \sum_{n=1}^{|\text{ACS}|} \beta_n \hat{x}_n \\ \sum_{n=1}^{|\text{ACS}|} \beta_n &= 1, \\ x \geq 0, y_s &\geq 0, \quad \forall s \\ \beta_n &\in \{0, 1\} \quad \forall n \end{aligned}$$



Bounds

Upper Bound

$Z_{HN}^*(I) \equiv$ The optimum objective function value to the ISP.

$Z_{HN}^*(C) \equiv$ The optimum function value of the LP relaxation.

$Z_{HN}(F_n) \equiv$ The objective function value of the n^{th} integer feasible solution

$$\text{where } Z_{HN}(F_n) = c\hat{x}_n + \sum_{s=1}^{|S|} p_s f_s y_s \text{ and } n = 1 \dots |ACS|$$

Let $Z_{HN}(F^*) = \text{Min} \{Z_{HN}(F_1), \dots, Z_{HN}(F_{|ACS|})\}$

Lower Bound

$$Z_{HN}^{LB} = \text{Max} \{Z_{WS}, Z_{HN}^*(C)\}$$

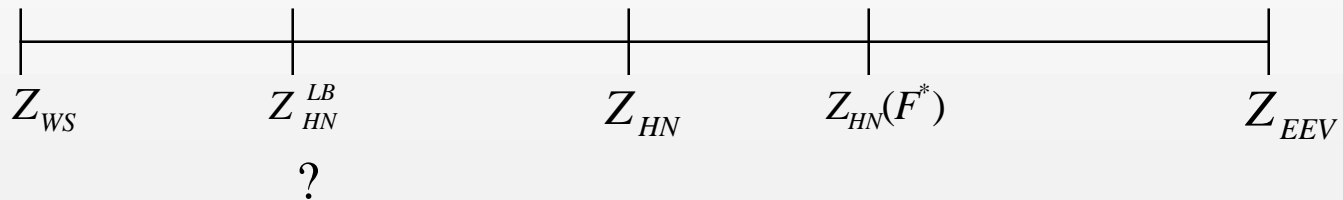
Then we have

$$Z_{HN}^{LB} \leq Z_{HN}^*(I) \leq Z_{HN}(F^*)$$



Qualifying the bound

For any 2 - stage SP minimisation model $Z_{WS} \leq Z_{HN} \leq Z_{EEV}$



The ratio $\frac{Z_{EEV} - Z_{HN}(F^*)}{Z_{HN}(F^*) - Z_{HN}^{LB}}$ measures the compactness of the bound

Computational Results

Model description

SCP-1 – Strategic supply chain, opening and closing of sites, DCs, lines.

SCP-2 – Trade-off between capacity acquisition and maintaining inventory.

Model Statistics

Model	Constraints	Variables (discrete/continuous)	Non-Zeroes
SCP-1	6768	2096/54400	1154034
SCP-2	574	216/240	1342





Network Dimensions		Planning Problem	
The number of Sites,	I :	8	
The types of packing line technology,	Y_C :	4	
The types of production line technology,	Y_R :	2	
The number of distribution centres ,	J :	15	
The types of DC line technology ,	Y_D :	2	
The number of Customer Zones ,	H :	30	
The number of Products ,	P :	13	
The number of time periods,	T :	6	
Model Statistics			
Logical Constraints: Sites, DCs opening and closing, Limit on number of Sites, DCs, and Lines		$m_1 = 968$	6768
Other Constraints: Production, Packing, Ordering, Transportation, Balance, Demand, and also Production and Packing Capacities.	Mixed	$m_2 = 850$	
	Continuous	$m_3 = 4950$	
Discrete Decision Variables: Sites, DCs, Production lines, Packing lines, DC lines.		$n_1 = 2096$	56496
Continuous Variables: Production, Packing, Ordering, Transportation, and Shortage quantities.		$n_2 = 54400$	
Non-zeros		1154034	
Scenarios		100	

Computational Results

Generation of integer solutions

Model	Number of scenarios	Number of Lagrangean iterations	Number of unique solutions
SCP-1	100	25	2050
SCP-2	43	25	143
SCP-2	43	15	85

Computational Results

Generation of integer solutions

Model	Z_{WS}	Z_{HN}^{LB}	$Z_{HN}(I)$	$Z_{HN}(F^*)$	Z_{EEV}	Quality
SCP-1	*	$2.52 \cdot 10^6$?	$2.64 \cdot 10^6$	$3.59 \cdot 10^6$	7.91
SCP-2(25)	$-1.23 \cdot 10^6$	$-1.23 \cdot 10^6$	$-1.21 \cdot 10^6$	$-1.18 \cdot 10^6$	∞	∞
SCP-2(15)	$-1.23 \cdot 10^6$	$-1.23 \cdot 10^6$	$-1.21 \cdot 10^6$	$-1.14 \cdot 10^6$	∞	∞

Comparison of computational time

Model	DEQ	Heuristic
SCP-1	158 hrs 35 mins	26 hrs 23 mins
SCP-2(25)	666s	980s
SCP-2(15)	666s	400s



Outstanding Issues

- Step-size in subgradient optimisation.
- The constraints to be relaxed.
- Parallelisation of the algorithm.
- Study the performance of the algorithm to SP test set.



Thank you !



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